

# Say What?

By Bonnie Baker

## Plans + DBRMS + Packages + Collections + Versions = Confusion

No matter how long programmers have worked with DB2 for z/OS and OS/390, they still ask me to tell them the difference between a plan and a package — and what in the heck a collection is. I planned to write a column on this topic so I could just point them to *DB2 Magazine*, Quarter 4, 2003 and let the column answer the questions. Well one column became two, and two morphed into three. And, if they keep changing DB2 (as they will do), the three may one day become four.

### THE EARLY DAYS

The complexity and breadth of this topic is best understood from an historical perspective. Long ago, when DB2 for MVS (now OS/390 and z/OS) was in its infancy, many companies had different machines for development and production. Development often took place on an older model machine, and the LOAD modules were moved onto a bigger production machine. In fact, this process still takes place in some companies if you take a broader perspective — some companies develop software offsite and then run it in entirely different locations for entirely different corporations. Come up with a need and companies will fill that need in unexpected ways.

COBOL was the language of choice in the old days; however, companies rarely had a COBOL license on the production machine. So, DB2 had to allow development on one box *with* a COBOL compiler but without a DB2 subsystem and allow production runs on a different box with a DB2 subsystem but *without* a COBOL compiler. That's how the concepts of the DB2 Precompiler and the database request module (DBRM) were born — out of necessity.

Programmers could write COBOL programs on the smaller, non-DB2 machine. They could embed SQL between the COBOL SQL delimiters (`EXEC SQL` and `END EXEC`) in both the COBOL working storage and procedural divisions. When they were ready to compile the program to look for COBOL errors, they'd run the program through the DB2 Precompiler to strip out the SQL, leaving only COBOL.

### THE PRECOMPILER

The DB2 Precompiler does *not* need DB2 to run. It carries out three primary tasks as it reads the program serially, top-to-bottom, looking for DB2 delimiters.

First, if the delimiters surrounded an `INCLUDE` statement, the Precompiler goes to the `INCLUDE` library named in the job control language data definition statement and pulls the included `MEMBERNAME` into the program. This *function* is the same as a COBOL `COPY MEMBERNAME`, but the timing is different. COBOL `COPYBOOKS` get copied in at `COMPILE` time; DB2 `INCLUDES` get copied in at precompile time. The only difference between an SQL `INCLUDE` and a COBOL `COPY` is timing. The most common item `INCLUDED` in a program was (and is) a `DCLGEN`. `DCLGENS` are structures that describe a table. One `DCLGEN` is usually

included for each table that the program will access at run time. Each `DCLGEN` is a two-part structure consisting of a `DECLARE TABLE` statement, which describes the table in DB2 SQL language, and a COBOL structure that describes the table using an 01-Level COBOL working storage structure (much like a typical copybook for a VSAM file).

Second, if the delimiters surround an SQL statement, the precompiler does a very basic syntax check to make sure that the column and table names are valid (that they're spelled correctly and that the columns and the table exist). Many DBAs and programmers think that this validation is done by reading the DB2 CATALOG, but they're wrong. Remember, the precompiler doesn't need DB2 or its CATALOG. DB2 might not even be installed on the machine. The DB2 Precompiler uses the top part of the `DCLGEN` to validate the SQL syntax.

The third, and most important, task performed by the DB2 Precompiler is to split the program into two parts: a COBOL and a DB2 part. All of the SQL that the programmer carefully embedded is stripped out of the program and put into its own partitioned data set (PDS) member, called a DBRM. A single program containing two languages, COBOL and SQL, goes into the DB2 Precompiler and two pieces come out. Twins, but fraternal twins — much like Arnold Schwarzenegger and Danny DeVito. Arnold looks just like his COBOL mother, and Danny looks just like his DB2 father. COBOL Arnold, with all of the SQL commented out, goes down one path in life. SQL Danny, containing only SQL, goes down a different path in life.

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The twins, separated at birth, have a tendency to lose each other. To help the twins find each other later in life (in other words, at run time), the precompiler engraves each with identical tattoos. The tattoo is carried forward with COBOL Arnold, through compile and link edit, into the `LOAD` module in the `LOAD` library. The tattoo is part of the run-time executable code of the `LOAD` module. The same tattoo is carried forward with SQL Danny through `BIND`. `BIND` is to SQL what `COMPILE` is to COBOL. The purpose of COBOL `COMPILE` is to come up with run-time code for the COBOL. The purpose of `BIND` is to come up with run-time executable code for the SQL. Both sets of code bear identical tattoos (timestamps or consistency tokens).

So, the COBOL twin becomes a transportable load module in the COBOL `LOADLIB` and the SQL becomes a transportable DBRM in the `DBRMLIB`. Just as the COBOL twin had to be compiled, the DBRM twin has to go through `BIND` to create the run-time executable code for the DB2 portion of the COBOL program and put that executable code into the "right" DB2 subsystem.

## ALL ABOUT BIND

`BIND` connects to the DB2 in which the program's `LOAD` module will run, reads the DBRM serially, and then performs three tasks.

The first of the `BIND` tasks is an authorization check. DB2 must make sure that the programmer has the `BIND` authority *and* the SQL authority to perform the requested SQL task (for example, updating the payroll master). When using standard authorization procedures, DB2 won't let you `BIND` a DBRM if you don't have the authority to execute the SQL that's in the DBRM. This is why you may have the authorization to `BIND` in development (accessing development tables) but don't have authorization to `BIND` in production, where the SQL accesses production tables. The second `BIND` task is a bit redundant. `BIND`, like precompile, must also check the syntax of the SQL, but the `BIND` check is more sophisticated. Instead of using the top, `DECLARE TABLE` portion of the `DCLGEN`, `BIND` uses the DB2 CATALOG table information to make sure that the column names are valid, that comparisons are numeric-to-numeric, and so on. This second syntax check occurs because you can't trust the one done by the precompiler because the precompiler check used the `DCLGEN`. You could have a `DCLGEN` and not have the DB2 table.

The third, and most important, `BIND` task is to come up with run-time instructions for the SQL in the

DBRM. Each SQL statement is parsed and all of the possible (realistic) methods for retrieving the desired columns and rows from the table are weighed, measured, and evaluated based on possible estimated I/O, CPU, and `SORT` overhead. A ton of information is used as input to the `BIND` process, not just `CATALOG` information put there by running the `RUNSTATS` utility. `BIND` input includes, for example:

- Indexes (what columns are in the indexes?)
- Columns (how long is this column and how much room will it occupy in a `SORT` record?)
- System resources (how big are the system resources, buffer pool, and `RIDPOOL`?)
- Processors (how big are they and how many engines do they have?)
- DB2 (what release is running?)
- Parameters (what are the values of the `BIND` parameters?)

After all that input (and more) is weighed and compared, the cheapest, most costeffective access path is chosen, and the runtime instructions for that one path are created. (Interestingly, DB2 `BIND` sometimes generates instructions for more than one path.) This process is called optimization, and it's repeated for each SQL statement in the DBRM until all access paths are decided and the run-time instructions are created for each. As the optimizer decides on each path, writes are done to DB2.

`BIND` checks to see if you bound with the parameter `EXPLAIN(YES)`; if so, it writes documentary evidence about the chosen path to the `PLAN_TABLE` and to the `DSN_STATEMNT_TABLE` for your edification.

`BIND` also writes a lot of information to multiple `CATALOG` tables, documenting the fact that the `BIND` did occur. In fact, the tattooed DBRM, which is not used at run time, is moved into the `CATALOG`. Objects chosen by the optimizer are documented in the `CATALOG` in cross-reference tables. And `BIND` parameters are recorded in the `CATALOG` also.

## WHERE ARE THE INSTRUCTIONS?

It's interesting that the actual instructions for the access path are not written to the `CATALOG`. You can't look at information in the DB2 `CATALOG` to figure out whether a query will do synchronous or asynchronous reads at run time. You can't tell if the query will match on three columns of an index or five. The actual run-time instructions aren't stored in the `CATALOG`. They're definitely not in the DBRM, which is input to the `BIND`. So, where are they stored?

This question is one of the many reasons that this column grew into two parts. Where in the heck are the run-time instructions? Should you ever use the `PLAN_TABLE` at run time? Are the run-time instructions in a package, a plan, or a version? All of the above? The infamous "it depends"? Where are the instructions stored while you wait for the `LOAD` module to run? How has this process changed over time? What program preparation options are currently available in version 7? What's coming in version 8?

Stay tuned for answers to these and other questions.

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**Bonnie Baker** is a consultant and educator specializing in applications performance issues on the DB2 OS/390 and z/OS platforms. She is an IBM DB2 Gold Consultant, a five-time winner of the IDUG Best Speaker award, and a member of the IDUG Speakers' Hall of Fame. She is best known for her series of seminars entitled "Things I Wish They'd Told Me Eight Years Ago" and the Programmers Only column in *DB2 Magazine*. You can reach her through Bonnie Baker Corp. at 813-876-0124 or

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## Say What? Part 2

By Bonnie Baker

### How the separated twins - modified source code Arnold and DB2 SQL Danny - find each other again.

In the last issue of *DB2 Magazine* I began writing about the program preparation process for DB2 Universal Database (UDB) for OS/390 and z/OS. If you haven't read Part 1, please do so now. Part 2 will make much more sense if you do. If you're still confused after reading Parts 1 and 2, maybe, just maybe, next quarter's Part 3 will clear up any remaining confusion and make you a program preparation expert.

By the end of Part 1, we had done a `PRECOMPILE` to separate our COBOL program into twins: Arnold (that's Governor Arnold now) as the modified source code (the SQL commented out and `DCLGENS` now `INCLUDED`) and Danny as the DBRM containing the SQL that used to be in the source code. The COBOL source code (without any SQL) was compiled into run-time executable instructions for the COBOL portion of the program; DBRM Danny (all the SQL that used to be in the source module) went through `BIND` to generate run-time instructions for the SQL in our COBOL program.

I didn't mention how Arnold, over in the COBOL `loadlib`, will find his long-lost twin — nor did I mention exactly what we were binding and where we would put it when we finished. And that's what this column is about.

#### The Modified Source Code (Arnold)

At precompile time, when the SQL was stripped out of our program and moved into the DBRM (Danny) leaving only COBOL in the modified source code, you must have wondered how COBOL Arnold would ever find DB2 SQL Danny; in other words, how the COBOL would ever execute any SQL.

The explanation is simple. All of the executable SQL (and not all SQL is executable — a `DECLARE CURSOR` isn't, for example) in the COBOL program was replaced with a `CALL` statement (in our example, a COBOL `CALL`). The modified source code, complete with its calls to DB2 SQL Danny, was compiled and "linked" into a `LOAD` module. When this `LOAD` module is executed and hits a paragraph that once contained SQL (but now contains a `CALL` to Danny), there will be run-time executable code that knows how to link to a COBOL-DB2 interface module and connect to DB2, where it will find Danny and the run-time executable code for the SQL statement that previously had been in the paragraph.

Remember that we tattooed the twins? Well, this `CALL` contains the information DB2 needs to confirm not only that this `LOAD MODULE` is Arnold (complete with requisite tattoo), but also that it's the exact same Arnold that came out of the exact same precompile step as Danny. The `CALL` looks at the `PLAN` named in the execute statement of the job control language (JCL) and searches for the Danny out in DB2 with the same tattoo.

#### The DBRM (Danny)

At `BIND` time, DB2 created run-time executable instructions for the SQL portion of the program. But where are those instructions, and what are they called now that the term DBRM no longer applies?

The truth is, you have a choice. You can `BIND` the instructions for the SQL that was in the DBRM into a `PLAN` (the old way), or you can `BIND` the instructions into a `PACKAGE` (the not-so-old-but-no-longer-new way). The reason for this choice is historical. Back when knights were bold, dragons walked the Earth, and DB2 and some of us were young, DBRMs were bound directly into `PLANS`. In today's DB2 (since V2R3), there are two ways of doing BINDs. You may continue to `BIND` DBRMs into `PLANS`, or you may `BIND` DBRMs into `PACKAGES`. With the second option, you keep your `PLAN` but use it only as a search chain. This column explains why things changed in V2R3. A future column will explain how things have continued to change in v.7 and v.8.

In the early releases of DB2, the DBRM (SQL originally embedded in our COBOL program but separated at precompile time into a PDS member) was bound into a `PLAN` (`PLANA`). This method worked just fine as long as the program was a standalone program. You coded the JCL to execute program `PGMA` naming the `PLAN` "`PLANA`," and at run time the twins found each other. However, things got a bit complicated when `PGMA` needed to `CALL` `PGMB`. Because only one `PLAN` could be named in an execute statement, the `PLAN` had to contain run-time instructions for both `PGMA` and `PGMB`. This problem was solved by having the `BIND` instruction for `PLANA` name a memberlist; the DBRMs for both `PGMA` and `PGMB` were listed as members. And if `PGMB` called `PGMC`, then the three would be listed as members. And if `C` called `D`, which could call `E`, `F`, `G`, or `H`, which could call `I`, `J`, `K`, `L`, or `M`, which could .... Well, you get the idea.

Memberlists got longer and longer. What were (and still are, if you cling to the old technique) the drawbacks of having a very long list?

1. Remember that DB2 authorization and SQL syntax are checked at `BIND` time. Access path alternatives are also weighed, the least-cost path is chosen (based on current statistics and system resources), and run-time instructions are created for the chosen path. Well, if the `PLAN` contains one member, `A`, this process should be quick. But what if there are 20 or 50 or 500 DBRMs in the memberlist? The `BIND` could take hours.
2. So, the `PLAN`, which took more than a while to `BIND`, contains 500 members. What if one of the programs, `PGMA`, changes? When the source code changes, the program must be precompiled. And precompile changes the tattoo and creates a new DBRM. That new DBRM must be bound into the `PLAN`. When the `PLAN` is bound, all 500 DBRMs (not just the modified `PGMA`) will be rebound. It could take hours to `BIND` a `PLAN`, even though 499 of the 500 programs haven't changed.
3. Also, remember that `BIND` is an opportunity to reassess and change the access paths for not only the modified program, but also every single program in the memberlist. If one program changes, every program in the list will go through `BIND`.
4. What if you want to modify `PGMB` to call a new program, `PGMZ`? You must not only precompile modified `B` and new `Z`, you also must add new `PGMZ` to the memberlist and `BIND` the whole list — all 501 DBRMs.
5. You want to remove `PGMT`? Edit the memberlist and then `BIND` the `PLAN` again with the remaining 500 members in the list and wait impatiently while DBRMs that haven't changed go through the `BIND` process.
6. Okay, modified program `A` turns your processor over on its back, casters up. You want to fall back to the original version of `PGMA`. And exactly how would you do that quickly? If (and it's a big if) you have the old DBRM with its old tattoo, you could move it into the `DBRMlib` and `BIND` the entire 500-member list (even though only `A` had regressed) and replace the new loadlib member `A` with the prior loadlib member `A`. Or, if you have the old source code for `A`, you could precompile it to recreate both the modified source code and the DBRM, and then `COMPILE`, `LINK`, and `BIND`, which would `BIND` all 500 DBRMs in the member list.
7. Program `Q` has to run against two sets of tables, one for Corporation 1 and a second for

Corporation 2. The sets of tables have identical names but different high-level qualifiers. You could use synonyms, but they're unwieldy; binding during a time when a synonym points to the wrong set of tables could cause disaster.

These (and other) quandaries faced many DBAs in the days before V2R3 and the advent of PACKAGES.

## Packages and Collections

In V2R3, the DB2 developers solved the seven problems I listed (and many more) by introducing another layer in the program preparation procedure. The precompile step still split the program into tattooed twins, and the treatment of the modified source code stayed the same. But the DBRM could now be bound either into a `PLAN` (the old way) or into a `PACKAGE`. Although the relationship between a DBRM and a `PLAN` was one-to-many, the relationship between a DBRM and a `PACKAGE` was always one-to-one. Most of the work of `BIND PLAN` was moved into `BIND PACKAGE`. Therefore, if `PGMA` changed, only `PACKAGE A` would have to be bound.

If only one DBRM could be bound into a `PACKAGE`, but `PGMA` could still `CALL PGMB`, then a structure was needed to gather all of the `PACKAGES` into a searchable list. This structure became a `packagelist`, rather than a `memberlist`, bound into a `PLAN`.

This tiny change solved many, but not all, of the problems inherent in the "memberlist of DBRMs bound into a `PLAN`" technique. To solve a few more problems, IBM introduced the concept of `COLLECTIONS`. A collection is simply a way of grouping packages into meaningful (for you) groups. You could `BIND` all of the packages that must run against Corporation 1's tables into one `COLLECTION` and all of the packages that must run against Corporation 2's tables into another. Or, you could use `COLLECTIONS` to separate programs for different application areas, such as payroll and inventory.

Another use might be to separate programs bound with `ISOLATION UR` from programs bound with `ISOLATION CS`. `COLLECTIONS` are simply high-level grouping names to designate that this group of packages share something, anything, in common.

So, now, with the introduction of `COLLECTIONS`, our `BIND PLAN` process (think search chain) can now include a `packagelist` of fully qualified package names such as `COLLPAYROLL.PGMA`, `COLLPAYROLL.PGMB`, and `COLLCOMMON.PGMX`. Or, if we choose, we can just substitute an asterisk for the program name and list (`COLLPAYROLL.*`, `COLLCOMMON.*`). Then any program bound into the `COLLECTION` will be accessible by the named `PLAN`.

## Where Are Packages and Plans?

When you `BIND` a DBRM into a `PLAN`, or `BIND` a `PACKAGE` into a `COLLECTION` and then the `COLLECTION(s)` into a `PLAN`, the information must be stored somewhere safe inside DB2 until it's needed at run time. These items (`PLANS` and `PACKAGES`) aren't stored in the DB2 Catalog. Rather, they're stored in the DB2 Directory. In fact, you can think of the Directory as the DB2 loadlib for the Danny portion of your program, complete with tattoo.

## Run Time

At run time, the load module starts up and eventually hits a paragraph containing a `CALL` to DB2. This `CALL` contains information such as a description of the tattoo, the content of your SQL host variables (now populated), the statement number, and so on. The `CALL` invokes the COBOL-DB2 interface program, which connects to DB2. And if the run-time code necessary to execute your SQL isn't currently resident inside DB2 (in the `EDMPOOL`), we go to the buffer pool (BP0) assigned to the DB2 Directory and

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look there. If we don't find Danny there, we go to VSAM to disk to look in the `COLLECTIONS` named in the `PLAN` for the `PACKAGE` with the same name and the same tattoo, also known as the consistency token or timestamp.

And if you don't find the twin anywhere in DB2 (not that this has ever happened to anyone reading this column), you get a -805 error. If you're still using the older technique of binding DBRMs directly into `PLANS` via a memberlist, then an unsuccessful search for Danny will result in a -818 error code.

## But What About Versions?

In my next column, I'll give you some ideas about the various ways of using the concept of the `COLLECTION`. I'll also write about `VERSIONS` and how they're used. Version 7 introduced some more useful variations that will carry us gracefully into v.8, and I'll cover them all next quarter.

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**Bonnie Baker** is a consultant and educator specializing in application performance issues on the DB2 OS/390 and z/OS platforms. She is an IBM DB2 Gold Consultant, a five-time winner of the IDUG Best Speaker award, and a member of the IDUG Speakers' Hall of Fame. She is best known for her series of seminars entitled "Things I Wish They'd Told Me Eight Years Ago" and the Programmers Only column in *DB2 Magazine*. You can reach her through Bonnie Baker Corp. at 813-876-0124 or

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# Say What? Part 3

By Bonnie Baker

## What to do when you find yourself in a bind.

Two columns ago, I wrote about the way we used to do program preparation for DB2 for z/OS and OS/390. Therefore, if you haven't read Part 1 of this series, please do so now ([Programmers Only, Quarter 4, 2003](#)). And while you're doing that, you should also read Part 2 ([Programmers Only, Quarter 1, 2004](#)). I promise that this column will make much more sense if you do. And if you're still confused after reading the two prior columns and this column, let me know via email, and I'll do a Part 4 just for you.

When we finished Part 1, we had done a `PRECOMPILE` to separate our COBOL program source code into twins. Twin 1 was Arnold (oops, Governor Arnold now), the modified source code (the SQL now commented out and `DCLGENs` now `INCLUDED`). Twin 2 was Danny, the DBRM containing the SQL that used to be in the source code. The COBOL Arnold source code (with all the SQL commented out and the executable SQL replaced by a COBOL `CALL` to DB2) was compiled into runtime executable instructions for the COBOL portion of the program and became a load module in the COBOL `loadlib`. DBRM Danny (all the SQL that used to be in the source module) went through `BIND` to generate runtime instructions for the SQL and became a `PLAN` in the DB2 Directory database.

By the end of Part 2, we had learned how this program preparation process changed in DB2 V2R3. We learned how Arnold over in the COBOL `loadlib` found his long-lost twin, Danny, at runtime. We also found out about the difference between a `PACKAGE` and a `PLAN`. And, we learned that `PACKAGES` are stored in `COLLECTIONS`. Finally, we learned where these things are stored when waiting for the lead program to be run.

In this column, I'll give you some ideas about the various ways of using the concept of the `COLLECTION`. I'll also write about `VERSIONS` and how they are used. DB2 v.7 introduced more useful variations that will carry us gracefully into v.8. I'll give you the high-level view.

### Collection Ideas

In Part 2, I said that a `COLLECTION` is simply a way of grouping `PACKAGES` into meaningful (for you) groups. You could use `COLLECTIONS` to separate programs for different application areas, such as payroll and inventory. Another use might be to separate programs bound with `ISOLATION UR` from programs bound with `ISOLATION CS`. `COLLECTIONS` are simply high-level grouping names to designate that this group of packages share something, anything, in common.

`COLLECTIONS` enable you to organize your `PACKAGES` into like-kind groups. In DB2's younger days, with multiple DBRMs being bound into `PLANS`, all the DBRMs in a single `PLAN` had to be bound with the same

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`BIND` parameters. However, today you can `BIND` each `PACKAGE` into a `COLLECTION` that has a customized set of `BIND` parameters associated with it. An example would be to `ISOLATE` all programs using `REOPT(VARS)` in one `COLLECTION` and all programs using `OPTHINT` in another. `DEGREE(ANY)` is another `BIND` parameter that you may want to be a bit more vigilant in monitoring. An easy way of keeping an eye on programs (or children) is to put them in a room together. In other words, `BIND` parameters are much more granular today than they were when DB2 was young.

With the advent of `BIND PACKAGE` and the one-to-one relationship of a program to a package, we were given the ability to name the high-level qualifier for the tables accessed by the program. Therefore, the `DBRM` for one program could be bound into two different `COLLECTIONS`. The `DBRM` for program `ABC123` could be bound into a `COLLECTION` called `colcorp01`, using `corp01` as the table high-level qualifier. The same `DBRM` could be bound into a `COLLECTION` called `colcorp02`, using `corp02` as the high-level qualifier. Or, you could `BIND` the same `DBRM` into `colstress` to run it against stress test tables and `BIND` it into `colstress` to run against regular test tables. Or, you could `BIND` a `DBRM` into a `COLLECTION` called `colur` to use when you access read-only decision support tables using `ISOLATION UR` and into a `COLLECTION` called `colcs` when you use active production data. There are dozens of examples. Just use your imagination.

So, at runtime, whichever approach you chose, you now have a `PACKAGE` with the exact same tattoo/timestamp/consistency-token in two different `COLLECTIONS`. How do you tell DB2 in which collection to search for Danny? Normally, DB2 would search through all of the `COLLECTIONS` in the named `PLAN`. But, if you want to search only one `COLLECTION`, you simply tell DB2 in your program. You can specify which `COLLECTION` to search by using an SQL statement called `SET CURRENT PACKAGESET`. `PACKAGESET` is simply a synonym for `COLLECTION`. Therefore, if you set the current `PACKAGESET` to `colcorp01`, you will access `corp01`'s tables. If you set the current `PACKAGESET` to `colcorp02`, you will access `corp02`'s tables. And the beauty of this is that you only have to maintain one program.

## Versions

In Part 1, I mentioned the need for rapid fallback. Suppose program A is changed and moved back into production. Before the program was changed it ran in 10 minutes and never bothered anyone. After the change, all the other programs running at the same time are experiencing -911 timeouts. How do you fall back gracefully and rapidly to the prior version of the program?

At precompile time you can specify a `VERSION ID`. If the `VERSION ID` is the same as the current version, `BIND` will overlay the `PACKAGE` in its `COLLECTION`. But, if the `VERSION ID` is different from the current `VERSION ID`, `BINDING` the `DBRM` will produce a new `PACKAGE` that won't overlay the prior `PACKAGE` for the program. You'll have two `PACKAGES` for the same program in the same `COLLECTION`. If you also move the current `LOAD` module with its old tattoo timestamp into a different `loadlib` (`COBOL.BACKUP`), the compile will not overlay it. Then, when you compile the modified source code, you'll have a `LOAD` module with the new tattoo/timestamp in the current `loadlib`. If you execute the new `LOAD` module, you'll find the new `PACKAGE`. If the system suffers, you can cancel the job and move the old `LOAD` module back into production simply by pointing to `COBOL.BACKUP`.

## Copies

There are too many other nuances and possibilities to mention; however, one feature that may be useful is the ability to copy a package from one collection to another. If the statistics on your table vary greatly from daytime to evening or beginning of the month to end of the month, you can `BIND` a `PACKAGE` in a `COLLECTION` called `colday` or `colbegin` when the statistics in the `CATALOG` are representative of your daytime or beginning of the month table. You can then `COPY` that `PACKAGE` into another `COLLECTION` called `colnight` or `colend` when the statistics in the `CATALOG` are representative of your nighttime or end of the month table. `COPY` does a `REBIND` and uses the `DBRM` in the `CATALOG` as its input. Therefore, the

tattoo/timestamp doesn't change. If you check the time of day or the day of month at the beginning of the program, you can SET CURRENT PACKAGESET to the appropriate COLLECTION for DB2 to search for Danny.

## Versions 7 and 8

In v.7, the developers introduced a new way of doing program preparation. The newer compilers (such as the Enterprise COBOL V3 compiler) let you skip the stand-alone PRECOMPILE. Why? This new compiler has an application program interface (API) to the DB2 precompiler. As a result, there is one less step in the program preparation process. The COMPILER is now more aware of your DB2 calls, and your SQL can take advantage of some of the host language features. For example, your DCLGENS can now have COBOL structures that look more like COBOL copybooks. The DCLGENS, once limited to 01 and 10-level fieldnames can now have field definitions that are more than two-levels deep. A 10-level trandate can be broken into 15-level year, editchar, month, editchar, day. Also, the DCLGEN fields can contain OCCURS clauses, which is much more programmer friendly.

This new capability carries us gracefully into the recently announced DB2 v.8 for z/OS. If you want to exploit the new multi-row fetch in v.8, you must fetch your rows into array structures. With the new compiler and its API, DCLGENS may now have the needed arrays.

## Got it?

For those programmers who are just beginning their work with mainframe DB2, the multiple layers and possibilities of DBRMs, PLANS, PACKAGES, COLLECTIONS, and VERSIONS can be confusing. The seemingly never-ending timestamp errors can be frustrating. I hope these three columns will clear up some of the questions about these DB2 entities and introduce possibilities never before considered.

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**Bonnie Baker** is a consultant and educator specializing in application performance issues on the DB2 OS/390 and z/OS platforms. She is an IBM DB2 Gold Consultant, a five-time winner of the IDUG Best Speaker award, and a member of the IDUG Speakers' Hall of Fame. She is best known for her series of seminars entitled "Things I Wish They'd Told Me Eight Years Ago" and the Programmers Only column in DB2 Magazine. You can reach her through Bonnie Baker Corp. at 813-876-0124 or

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